

Quantifiers in Limits

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Quantifiers in limits

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Robinson's limit definition

The nonstandard approach to calculus eliminates two quantifier blocks in the limit definition. The standard definition of

$$\lim_{z \rightarrow \infty} F(z) = \infty$$

requires three quantifier blocks,

$$\forall x \exists y \forall z [y \leq z \Rightarrow x \leq F(z)].$$

A. Robinson 1960: For standard functions F , this is equivalent to the universal sentence

$$\forall x [I(x) \Rightarrow I(F(x))]$$

where $I(x)$ means x is infinite.

Quantifier hierarchies

Fix an ordered structure $\mathcal{M} = (M, \leq, \dots)$ with no greatest element.

Hierarchy of sentences in $L(\mathcal{M}) \cup \{F\}$:

Π_n : n quantifier blocks starting with \forall

Σ_n : n quantifier blocks starting with \exists

Δ_n : Both Π_n and Σ_n

B_n : Boolean in Π_n .

$$\Delta_1 \subset \Pi_1 \subset B_1 \subset \Delta_2 \subset \Pi_2 \subset B_2 \subset \Delta_3 \subset \Pi_3,$$

$$\Delta_1 \subset \Sigma_1 \subset B_1 \subset \Delta_2 \subset \Sigma_2 \subset B_2 \subset \Delta_3 \subset \Sigma_3.$$

Problem. Given \mathcal{M} , locate

$$LIM = \{(\mathcal{M}, F) : \lim_{z \rightarrow \infty} F(z) = \infty\}$$

in the quantifier hierarchy.

For each \mathcal{M} , it's Π_3 or lower.

Cases with low quantifier level

Theorem. *For every \mathcal{M} , LIM is not B_1 .*

Theorem. *If \mathcal{M} has universe \mathbb{R} and a symbol for each function definable in $(\mathbb{R}, \leq, +, \cdot, \mathbb{N})$, LIM is Δ_2 .*

Proof: \mathcal{M} has a symbol for a function $G(x, n)$ such that $\mathbb{R}^{\mathbb{N}} = \{G(x, \cdot) : x \in \mathbb{R}\}$.

LIM is equivalent to

$$\exists x \forall n \forall y [G(x, n) \leq y \Rightarrow n \leq F(y)].$$

The negation of LIM is equivalent to

$$\exists m \exists x \forall n [n \leq G(x, n) \wedge F(G(x, n)) \leq m].$$

Cases with maximum quantifier level

Theorem. *If \mathcal{M} is countable, then LIM is not Σ_3 .*

Theorem. *If $\mathcal{M} = (\mathbb{R}, \leq, \mathcal{N})$ with $\mathcal{N} = (\mathbb{N}, \dots)$, then LIM is not Σ_3 .*

Theorem. *If \mathcal{M} is saturated (or special), then LIM is not Σ_3 .*

K. Sullivan, Ph.D. Thesis 1974, showed that LIM is not Π_2 and not Σ_2 when \mathcal{M} saturated.

Theorem. *If $\mathcal{M} = (\mathcal{K}, I)$, \mathcal{K} saturated and $I = \{x : x \text{ is infinite}\}$, LIM is not Σ_3 .*

So Robinson's result for standard functions does not extend to arbitrary functions.

Infinitely long sentences

Given a set of sentences Q ,

$$\bigwedge Q = \{\bigwedge_n \theta_n : \theta_n \in Q\}$$

$$\bigvee Q = \{\bigvee_n \theta_n : \theta_n \in Q\}.$$

If \mathcal{M} has universe \mathbb{R} and a constant for each $n \in \mathbb{N}$, then LIM is $\bigwedge \bigvee \Pi_1$,

$$\bigwedge_m \bigvee_n \forall z [n \leq z \Rightarrow m \leq F(z)],$$

and LIM is $\bigwedge \Sigma_2$,

$$\bigwedge_m \exists y \forall z [y \leq z \Rightarrow m \leq F(z)].$$

Theorem. *If \mathcal{M} has universe \mathbb{R} , LIM is not $\bigvee \bigwedge B_1$.*

o-minimal structures

\mathcal{M} is o-minimal if every set definable in \mathcal{M} with parameters is a finite \cup of intervals and points. (Van den Dries 1984, Pillay and Steinhorn 1986).

Examples of o-minimal structures:

$(\mathbb{R}, \leq, +, \cdot)$ (Tarski 1939).

$(\mathbb{R}, \leq, +, \cdot, \exp)$ (Wilkie 1991).

Above plus restricted analytic functions
(van den Dries and C. Miller 1994).

Theorem. *If \mathcal{M} is an o-minimal expansion of $(\mathbb{R}, \leq, +, \cdot)$, LIM is not $\wedge \Pi_2$ and not $\vee B_2$.*

Proof uses recent results of H. Friedman and C. Miller (2005) on fast sequences.

Conjecture. *For every o-minimal expansion \mathcal{M} of $(\mathbb{R}, \leq, +, \cdot)$, LIM is not Σ_3 .*

Summary

Quantifier Level of LIM Over \mathcal{M}

$$\Delta_1 \subset \Pi_1 \subset B_1 \subset \Delta_2 \subset \Pi_2 \subset B_2 \subset \Delta_3 \subset \Pi_3$$

$\leq \Pi_3$	always
$\leq \wedge \Sigma_2$	$(\mathbb{R}, \leq, 0, 1, 2, \dots)$
$\leq \wedge \vee \Pi_1$	$(\mathbb{R}, \leq, 0, 1, 2, \dots)$
$> \Delta_3$	countable
$> \Delta_3$	(\mathcal{M}, I) with \mathcal{M} saturated
$> \Delta_3$	$(\mathbb{R}, \leq, (\mathbb{N}, \dots))$
$> \vee B_2$	o-minimal $(\mathbb{R}, \leq, +, \cdot, \dots)$
$> \wedge \Pi_2$	o-minimal $(\mathbb{R}, \leq, +, \cdot, \dots)$
$> \vee \wedge B_1$	$(\mathbb{R}, \leq, \dots)$
Δ_2	$(\mathbb{R}, \leq, +, \cdot, \mathbb{N}, \text{definable})$
$> B_1$	always

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